

Performance Analysis of Attribute Signature Management Framework(ASMF) for Digital is Document Security and Integrity

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Abstract— Data security and protection is the major issues Smart and robust authentication systems are need of today's scenario. Authentication and cryptosystems are some of the approaches which deal with user's identity verification and data confidentiality. Well known implementation area of hashing is digital signatures which deal with the documents originality and creators identity. It is modified with adding additional information related to the attributes of users and comes under attribute based signatures (ABS). We have made in depth analysis of some previous approaches in work [16] that were offering the ABS phenomenon and found some of the unsolved issues solution to which reduces the complexity, consumption and overhead associated with signature generations and verification. The work[16], we proposed a novel attribute signature management framework (ASMF) for handling the process of digital signatures with attributes and its predicate logic. The solution contains the versatile primitives to allow the signing parties having complete controls over the signature generation and the fine grained access security controls. It has a predicate logic attested with message and hides the signatures privacy. In this work we have collected stastical data that can justify the advantages of the proposed ASMF in work [16]. We have collected significant data in tabular form which can easily estimate the benefits of the approach. The basic purpose of this work is to analyse the performance of the framework proposed in paper[16], After analyzing the approach analytically and stastically some of the benefits were found which can be further enhanced and used in significant areas of industry and research in near future.

Keywords— Data Security, Digital Signature, RSA, MD5, Certificate X.509, Attribute Signature Management Framework (ASMF), Computational Complexity, Resource Consumption, Overhead;

I. INTRODUCTION

Cryptosystem deals with data security and protection which were designed in temporary fashion where the attacked systems are handled. All the developed models are mathematical strong to handle the unauthorized access and changes tried by the attackers. This mathematical problem is computational hard and the complexity of system needs to be assumed which was in traceable and used as a major building block for security system. Cryptosystems are further pored into symmetric and asymmetric nature depends on the key used. Some of the examples of such cryptosystems are DES, AES and RSA. Here the RSA approach is a public key cryptosystem. The tools developed with these cryptosystems are offering through random oracles for proving the security of a scheme. A random oracle is an oracle that when queried responds with a

random reply, subject to the condition that the reply is different for various queries but the same when the same input is queried again. Such an oracle has the desired properties such as preimage resistance and collision resistance. Signatures have existed as a means of authentication for ages.

It takes different formats according to the application. Until that point signing a document meant identifying the person. Signatures have been very useful in creating certified documents, but it has generally been necessary to rely on external evidence that the persons signing the document are qualified to do so. This is today's challenge. In the early days of digital signatures, signers owned a pair of keys, a public key known to everyone used as an input for the verification process and a private key known by the signer and used in the signature process. Then the scheme is proved secure under that random oracle assumption. Finally when it comes to actually implementing the scheme the random oracle is replaced with a strong hash function. Researchers believe that proving a cryptosystem secure under the random oracle is equivalent to proving the security of the scheme dependent on exploiting the hash function used. Research in Attribute based signatures followed a certain trend[16].

Cryptographers would identify an application that requires attribute signatures where that application would require certain properties that no existing cryptosystem provides. They would propose a totally new scheme with new security notions that serve such properties. The result was several schemes each serving a very specific application making it hard to employ in any other[16].

Understanding Digital Signature

If a user wants to send a signed document to another user, public key is used to identify the users and it was known to everyone and the secret key is not known. Using the secret key and a message, first user can create a signature and send it to another. The recipient user can verify the signature using the message and the public key. The signature contains several elements that look random to recipient. Among these elements is one which is referred to as a "Fingerprint". Second user (Recipient) knows the verifying procedure which will enable him to make use of the public key and the signature in order to recalculate the "Fingerprint". If what he calculated is equivalent to what he

got from first user (Sender) then accept signature otherwise reject it. A Digital Signature Scheme is a set of algorithms. To define the scheme we explain the algorithms[16].

- **Setup (k):** This algorithm creates two keys using a security parameter k . The first is the public key P_k known to all and the second is a secret key S_k given to the signer only.
- **Sign (M, S_k):** This algorithm is run by the signer (first user). He signs the message M using a secret key S_k and outputs a signature "Sig".
- **Verify (Sig, P_k , M):** This algorithm is run by the verifier (Recipient). He uses and the public key P_k to run the verification algorithm that will output either accept or reject.

Attribute-Based Signature (ABS)

A digital signature is the mathematical construction scheme for generating the authenticity information for verifying the digital documents or identity of users. A signature based document will prove that it was created and transmitted without any modifications and having proper privileges assigned to the sender. It uses the attributes from the set of user's information instead of any single feature for showing the signers identity. A valid ABS is the signature which verifies the message content, its predicate and users identity associated with it. It highlights the word single in this familiar security guarantee ABS signatures, as in most attribute-based systems, require that colluding parties not be able to pool their attributes together. Furthermore, attribute signatures do not reveal more than the claim being made regarding the attributes, even in the presence of other signatures. An attribute-based signature assures the verifier that a signer, whose set of attributes satisfies a (possibly) complex predicate, has endorsed the message. It is further divided into three types:

- **Group Signatures**

Group signatures are digital signatures that allow any member of a group to sign anonymously on behalf of the group and in case of a dispute, a trusted group manager can revoke that anonymity. Important concern to a group signature scheme is a group manager, who is responsible of adding group members and has the ability to disclose the original signer in the occurrence of ambiguity[16].

- **Ring Signatures**

A ring signature is similar in concept with group signatures but differs in three key ways. First of all, there is no way to revoke the anonymity of an individual signature (i.e. no one can tell the signer of a message not even the group manager). The next distinction is other group of users can be considered as a group without additional setup. The last distinction is that every user has a public and private key[16].

- **Mesh Signatures**

The idea can be considered as an addition to ring signatures, but with added modularity and a much comfortable primitives for expressing signer ambiguity. Intuitively, mesh signatures (as in ring signatures) need to be anonymous and unforgeable. The access structure can be

satisfied using different combinations of atomic signatures, once created the mesh signature will not release what particular subset was used[16].

Properties of ABS

An ABS scheme allows a verifier to decide on the set of attributes (s) he would like the signer to possess. The verifier sends the request to a group of possible signers as a monotone Boolean expression. Any member with sufficient attributes can sign. The scheme maintains certain properties as follows[16]:

- *No previous knowledge assumption*
- *Unforgeable*
- *Anonymous Identities*
- *Unlinkable*
- *Traceable*
- *Anonymous Attributes*
- *Coalition Resistant*
- *Separability*

II. RELATED STUDY

During the last few years the security of digital document is considered as major issues. Its protection is implemented using digital signature. We have gone through a variety of signature algorithms and attribute based signature is one of its recent approach. Here with this literature review we put a light on various existing approaches and their functional behavior to develop a better approach[16].

In the paper [6] attribute based cryptography is discussed for getting the fine grained access control. It holds and verifies the secret information associated with the digital documents. Also a great use of attribute based signature is shows here to develop the practical solution satisfying the predicate conditions. The problem associated with the authentication and anonymity in distributed access control system, dealing with resources consumption is very much important. The detected size of signature generated by existing ABS schemes grows linearly with the number of attributes used for forming a signing predicate. The paper also proposes the first two attribute-based signature schemes with constant size signatures. Their security is proven in the selective-predicate and adaptive-message setting, in the standard model, under chosen message attacks, with respect to some algorithmic assumptions related to bilinear groups.

Some of the paper shows the evolution of ABS based on the size of signature and the predicate logics as mentioned in the paper [7]. The length of signature mostly depends on the largest attribute which can be predicted by some attackers. Thus the paper proposes an attribute-based signature scheme with uniform size irrespective of the nature of attribute size. This scheme is unforgeable conditionally depends on the success probability of any polynomial time adversaries. It is proved to be unforgeable and unconditionally anonymous in the standard model. The security of our schemes has been proven under the standard computational Diffie-Hellman assumption.

In the paper [8] cryptographic methods are studied against the system scalability factors, access control and other key generation activities. Mainly it applies the

attribute based encryption for performing the storage over cloud servers. The approach of ABE is further divided in KP-ABE and CP-ABE. For preserving the security and data protection CP-ABE is more appropriate. To implement this the CP-ABE is applied to the data by service providers on the documents uploaded by the users. For implementing the signature approaches the ABE is converted to attribute based Signature (ABS) schemes. The paper also discusses the variants of ABS like Ring ABS, Group ABS and Mesh ABS. Thus, ABS proves to verify that the signer holds a subset of attributes satisfying that signing policy. It can be efficiently used in real scenarios like data sharing in cloud computing for certification and confidentiality with satisfying signing policies.

The paper [9] further works on specific elements of ABS implementation using designated signature verifier entity. The paper proposes a threshold attribute based signature verifier scheme and works as an effective model for assuring the identity of the sender. The scheme named as t-ABUDVS scheme consists of the following algorithms: t-ABUDVS= (Setup, Extract, PS, PV, DS, DV, Sim). The security of suggested scheme will be reduced to the hard problem in which the signature is constructed. The paper also put lights on the definition of the bilinear Diffie-Hellman problem.

The paper [10] presents a novel attribute-based signature scheme for recovering the corrupted or lost messages. While comparing the scheme with traditional approaches, this ABS scheme with message recovery is not having any requirements regarding the retransmission of the original copy of message for verifying the signatures validity because of its nature by which the original message recovery from signature is prevented. The scheme also reduces the total length of the original message with appended signature. The paper deals with its three contributions i.e. attribute-based signature with message recovery, bilinear pairing construction and signature applicability for larger sized messages. It also support existible threshold predicates and are proven to be existentially unforgeable against adaptively chosen message attacks in the random oracle model under the assumption that the Computational Die-Hellman problem is hard.

In the paper [11] a new ABS technique based on the IRMA attribute-based authentication is proposed for healthcare industry. The proposed scheme of IRMA has an smart card based effective and practical implementation prototypes working successfully. The paper extends the existing functionality along with new implementation of ABS for IRMA devices. It also give a study on practical issues that arise due to the introduction of the signature functionality to an existing attribute-based authentication scheme, and we propose possible cryptographic and infrastructural solutions. For implementing the solution the paper gives a design evaluation using its use cases analysis.

While discussing the ABS most of the authors claims to deals with unforgeability which deals with signing key. A user can only be able to put its signature if he is having a complete control of its key and policy. It is having different entities like signature authority, verifier, users, receiver. The work is made clear for applying the signature

on group of users with leaking secrets and is having expressive predicate. It applies the policy of perfect privacy anonymity for hiding the users and its signature identity. Some of its extended variants are Traceable ABS (TABS), Decentralized TABS [12]. The work is also presented with practical implementation and efficient construction of suggested approach. It also prevents the expressiveness of designed policies with static assumption.

III. PROBLEM IDENTIFICATION

Internet technology is changing very rapidly with evolutionary architecture on which the webs that are achieving the practical realities. The Internet of Things (IoT) [13] is one of such area which deals with the concept of day to day electronic devices monitoring device, sensors, and home appliances accessing the internet for their effective use. It works towards transforming the devices to smart working objects which was flexible in nature and analyses process which was there with the internet. As the technology is going the requirement of security is also generating the demanding situations. Existing security primitives is not applied directly on IoT due to their segregated and heterogeneous standards and communication stacks. Moreover, the high number of interconnected devices arises scalability issues; therefore a exible infrastructure is needed able to deal with security threats in such a dynamic environment [16].

While the signature is one of such process which is necessary to assure the users identity and its document confidentiality. Considering the attribute based signature (ABS) schemes, after analyzing the approaches and the material we have found some of the current issues which needs to be resolved for implementing an more robust security controls over IoT. These are[16]:

- (i) High computational cost related to the signature process and which is directly proportional to the predicate formula which will not work for hand held devices. In some cases the signature predicate formula is smaller but the document confidentiality is more thus here the security is compromised with existing scheme [15]. Thus to reduce overhead is the primary task with this work.
- (ii) Some of the approaches will generate the variable sized signature but the signature for the user will be same sized and it depends on the attributes. Thus the signature must be designed in such a way which will generate the constant sized signatures to reduce the overhead and its associated complexity.
- (iii) The existing schemes sometimes expose the key, this must be improved to gain more robustness against the unforgeability and attribute signature privacy issues. This can be handled by introducing the random function for key generation and handling using key exchange mechanism.

After analyzing the problems with existing mechanism a new attribute based signature (ABS) is required with better construction privacy, lesser computational cost and reduced signature size with same or higher protection levels[16].

IV. PROPOSED WORK

Attribute Based Signature (ABS) offers the benefits to the user by passing the selected attributes from set of attributes for generating the signature with higher strength. It offers the phenomenon of anonymity for serving the privacy of user’s identity and the signature. Here the key revocation is associated with attribute generation which will improve the performance of ABS. But it was a challenging situation for the signature verifier because he doesn’t know the users selection towards its attributes and hence the verification process is not directly applied here. This paper proposes a novel attribute signature management framework (ASMF) which deals with all the above mentioned problems along with its efficient implementation. It shows an improvement over the key revocation using a designated authority for dealing with users selected attributes. The designed authority works as interacting medium between the verifier and the user and applies a negotiation of attributes. For generating the signature the user ask the intermediate for its secret key share along with a revocation check for its identity and the desired attributes. Here the key exchanges are handled using the RSA algorithms[16].

The process starts with forming a users group interested in making the digital signature based documents.

All the users are having different set of properties associated with their identities and system usages habits. We called them as attributes. These attributes are extracted from the users and stored to attribute store. Later on the data fragment passed by the user on which the signature has to be attached are passed to the hash generation modules which is having an MD5 digest algorithm working in parallel with the signature predicate logic module. It contains the additional information about the users in the form of their attribute passed for documents identity verification in the form of signature[16].

The hash algorithm calculates the digest using this predicate logic. Now the RSA cryptosystem is used for encrypting the generated digest using the signer’s private key. This private key will prove the authenticity of users along with its attributes. After applying the encryption X.509 certificate is added to hold the authentication information of users. This authentication information verifies by authentication server as identity check. All the temporary data generated by the system is stored with signature in the data repository. It is used to recreate the signature if repeat demands generated by the user for other document. It also holds the predicate logic decided for signature generation. Now the signed message of document is transmitted over the open channel[16].

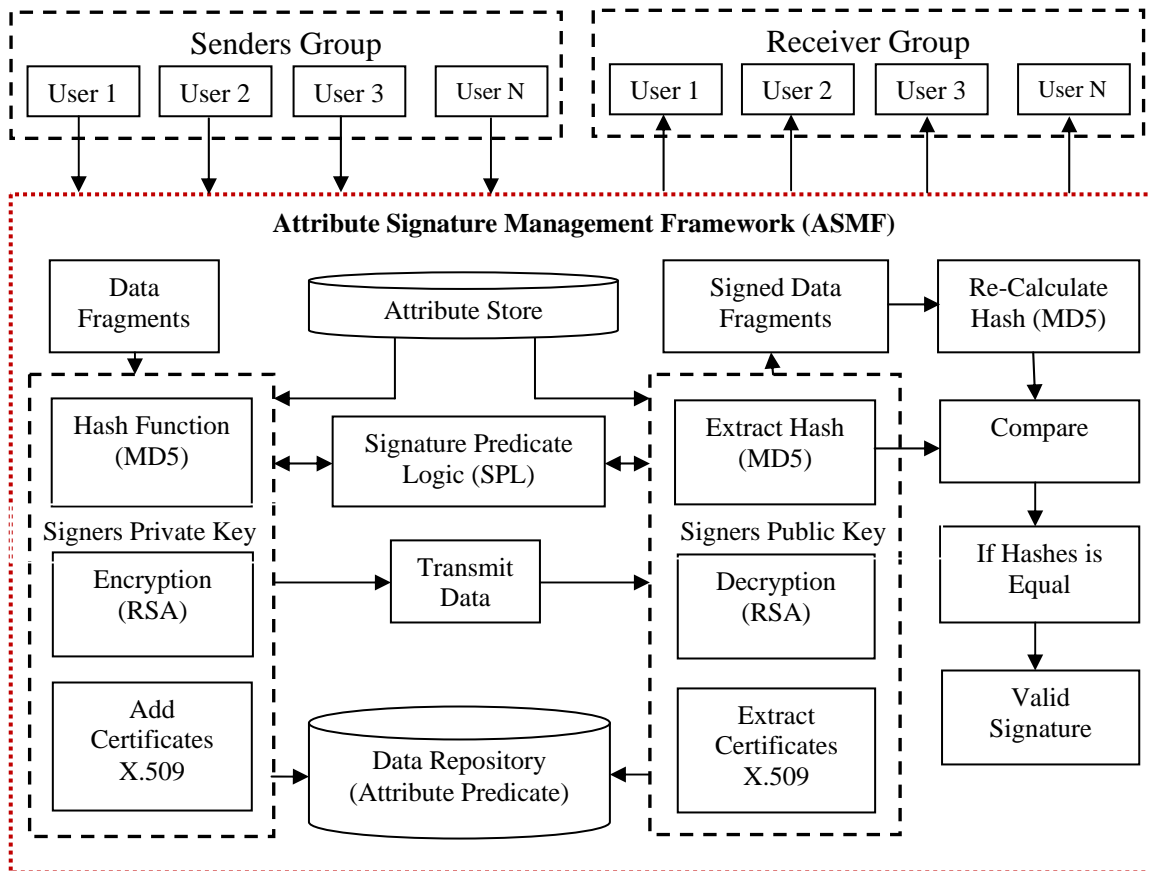


Figure 1: Proposed Framework for Attribute Signature based Management.

Now the second phase of proposed ASMF framework is for verification process. This phase starts with extraction of digest associated with the data fragments. Before extraction we need to decrypt the digest using signer's private key maintained with the public key register. Once the digest is decrypted then the associated hash with the data is extracted and the original data is passed again for recalculating the hash using MD5 algorithm.

The recalculated hash is compared with the extracted hash along with the user's attributes verification. If the attributes are matched with the user's identity associated stored then the validity of the signature and the sender is confirmed. Finally the message is made available to the receiver with its proof mentioned with the certificate.

V. RESULT ANALYSIS

Algorithm developed by us and implemented successfully. The results are shown for digest, certificates, encryption & decryption for text as well as for different type of file. We have separated the message as text and file for encryption. The aim with ASMF is towards providing more security and robustness against the traditional breaches. These extensions are kept in mind at the time of implementing the suggested concept. Now, it's the time to show how better the approach can work while comparing with their competitors. At the evaluation point of view the approach seems to satisfy all the constraints of data isolation, confidentiality, integrity and authentication. If the system provides effective security control but let's open the trapdoors for the malicious user then the confidentiality of the system can be loosed. It has a risk associated with outsourced environment. Different type of users has different type of permissions to access the information and functionalities in the application. Because of that reason, the system has full security for its data and functionality.

We have tried to get significant statistics data from the application we have developed. We need some relevant data to justify the efficiency of the work. We are going to get the statistics by executing the application in different cases. We will try to get the data from the application for different users, files, size, platform etc. so that we can have a variety of statistics to study the benefits of the proposed approach, that will off course act as user attribute record as shown in the table below.

Table I: User Attributes Record

S. No	User Name (Attribute Value 1)	File Name (Attribute Value 2)	Size Attribute Value 3 Bytes	System Name (Attribute Value 4)	Platform (Attribute Value 5)	IP (Attribute Value 6)
1	User1	User_File_Logs.txt	292	SATYADEEP	Windows 7 or Above	127.0.0.1
2	User2	Administration.docx	14715	SATYADEEP		127.0.15.6
3	User3	licence_4.jpg	38002	SATYADEEP		127.63.23.54
4	User4	Publication (Main)-1.docx	124871	SATYADEEP		127.45.68.95
5	User5	Data Mining.docx	39310	SATYADEEP		127.86.32.47

Table Interpretation: Above table-I is showing the different attributes used in the algorithm to apply the proposed solution. In this table we have user different users and the users uses the application for the variations of the

file according to its type and size. We have also recorded the other useful attributes for ASMF like system attributes, platform, the IP address etc. This table is the initial input the ASMF as the attributes shown above in the table will be processed further for the hash calculation, encryption and digital signature as the phases of the algorithm.

Table II: Secure Transmission Statics

S. No	File			Time (ms)					Overall
	Name	Size (bytes)	ID	Hash	Public Key	Private Key	Encryption	File Sending	
1	Data Mining.docx	39310	1045	34.45	2.08	1.063	2402.15	2298.02	2440.96
2	3 rd sem syllabus.pdf	15982	1046	36.99	2.25	1.181	5259.56	2190.81	5300.81
3	Paper856 95-700.pdf	173616	1047	0.77	1.29	1.062	2254.59	2033.66	2258.20
4	Networki ngquestio ns.docx	37038	1048	0.38	1.54	1.144	2034.08	1660.55	2037.07
5	Paper856 95-700.pdf	173616	1049	0.75	1.53	1.052	2396.67	1631.66	2400.06

Table Interpretation: In this table II we have collected the core statistics found by applying the ASMF through the application. This table is showing how much time(in ms) the proposed algorithm took when gone through different phases along with the overall time duration of the process. In this table we have applied ASMF on the users and their attributes mentioned in table I.

The proposed algorithm ASMF is constituents of steps as mentioned the proposed architecture of the approach like hash calculation, key generation, RSA encryption and final upload, the table above is clearly showing the variation of the time duration that taken by the proposed ASMF on the variety of files and sizes. The time duration we recorded for different heads as mentioned are in mil. seconds(ms)

Table III: Secure Extraction Statics

S. No	File			Extraction Time (in ms)				Overall
	Name	Size (bytes)	ID	Hash Extract	Public Key Check	Decryption	File Download	
1	Data Mining.docx	3931	1045	7.158	9.384	2.283	9.514	29.587
2	3 rd sem syllabus.pdf	15982	1034	36.110	111.386	91.785	48.683	221.769
3	4.jpg	897738	1028	19.650	17.650	2.862	38.277	80.960
4	Publication .docx	124871	1040	18.999	24.571	5.985	77.236	129.906
5	Assignment 2.pdf	120233	1032	20.285	24.171	6.164	64.118	118.105

Table Interpretation: In this table we are trying to depict the extraction time of the ASMF algorithm. In this approach we have proposed the data integrity and validity mechanism using Hash, RSA and digital signatures. According to the approach, to ensure integrity we need to recalculate the Hash and verify it with the previously calculated hash. Then for the next step we will the key for ensuring authenticity of the file. The table above is showing time to recalculate the hash and decrypt the file to regain the original form of file. The table above is showing the verities of file and sizes as shown in previous tables.

Table IV: Qualitative Analysis Comparison with Existing Approach

S. No	Approach Name	Feature Difference	Time	Size	Complexity	Efficiency
1	MD5 (Existing)	Single	High	Variable	High	High
2	RSA (Existing)	Single	Optimal	Variable	Low	High
3	Digital Signature (Existing)	Single	Low	Fixed	Low	Low
4	Certificate Generation (Existing)	Single	Optimal	Fixed	High	Low
5	Proposed ASMF	Multiple	Optimal	Variable	Low	High

Table Interpretation:

The above table shows the qualitative analysis of proposed ASMF algorithm along with the existing individual approaches. The table shows the behavior of the system during the execution for different cases of file, types, sizes and performances. By the table it is clearly shown that the suggested approach is setting mark for future research in the area of security. As of now the Hybrid approaches is not robustly analyzed but with some more cases and equivalency testing we could able to get the accurate analysis.

VI. IDENTIFIED BENEFITS

At the initial level of our research we get the following benefits.

- Our technique provides a feasible way to realize the “piecewise key generation.
- To allow for high efficiency and flexibility.
- Small number of exponential computations is enough for user signing.
- Less Complexity
- Efficient Data Outsourcing.
- Computational Cost is low.
- Better performance in server end, so that the data outsourcing becomes easier in data transferring communities.
- It retains reliability on third party location
- Client ensures about data storage in safe manner and unauthorized access.
- It protect data from different attacks at client end
- It might be became innovative approach at client end in cloud platform for different application domains.
- The storage & computation cost can be minimized.

VII. CONCLUSION

Data security depends upon the handling mechanism used to exchange the information between different ends. For verifying the authenticity of digital documents we need to verify the senders and the document both after receiving it. The file deals with it Is digital signature. Providing the user with more facility towards the data security and a assurance a new field is working with exiting mechanism known as attribute based signature. We had made a study on various attribute based signature mechanism and analyze their working capabilities to detect some of the unsolved issues. Mainly the computational complexity associated with these algorithms is very high and the types of resources they are consuming is also high. Thus a new approach is required to solve the issues. In this paper we

propose a novel attribute signature management framework (ASMF) to overcome these issues. At the analytical level of evaluation we are getting the effective outcomes which could be latter verified by its prototypic implementation developed in near future.

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